

Name: \_\_\_\_\_

Be sure to show your work!

**1. (25 points)** “Fun” with Truth Tables!

- (a) Fill out an abbreviated truth table for the following statement:  $(a \vee \neg a) \rightarrow (b \wedge \neg b)$ . Circle your concluding truth values and the correct answer:

This statement is a    **Contingency**    /    **Contradiction**    /    **Tautology**.

- (b) Show that  $\neg a \vee \neg b$  and  $\neg(a \wedge b)$  are logically equivalent.

- (c) Consider the following statement:  $\neg a, a \rightarrow b \vdash \neg b$ . Circle the correct answers.

**IS**    /    **IS NOT**    a theorem of  $L$  by the    **Soundness**    /    **Completeness**    theorem.

## 2. (25 points) System L

(a) I have provided a proof of Lemma L6 below – except the justifications are missing.

Fill in the justifications and please **be specific!**

**Lemma L6**  $A \rightarrow (B \rightarrow C), B \vdash_L A \rightarrow C$

1:  $A \rightarrow (B \rightarrow C)$

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2:  $B$

---

3:  $A \rightarrow B$

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4:  $(A \rightarrow (B \rightarrow C)) \rightarrow ((A \rightarrow B) \rightarrow (A \rightarrow C))$

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5:  $(A \rightarrow B) \rightarrow (A \rightarrow C)$

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6:  $A \rightarrow C$

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(b) Prove Lemma L10:  $\vdash_L (\neg B \rightarrow \neg A) \rightarrow (A \rightarrow B)$ .

[You may use the deduction theorem if you wish.]

*Hint/Suggestion:* Axiom 3 should be helpful.

**3. (25 points)** Models, Variables, Free for...

(a) Consider the formula:  $\forall x (A(f(\underline{a}, x), z) \rightarrow \exists y (B(y) \wedge C(x, y, z)))$ .

- i. Underline the scope of each quantifier in the above formula.
- ii. **Circle** all of the **bound variables** in the above formula.
- iii. Circle the correct answers and fill in the blanks below:

The formula above **IS** / **IS NOT** a sentence since \_\_\_\_\_.

The term  $t = g(y, \underline{a})$  **IS** / **IS NOT** free for  $z$  in the formula above.

(b) Consider the formula:  $(\forall x A(x, \underline{c})) \rightarrow (\exists y B(y))$ .

Translate the above formula into *plain English* when our model consists of the universe of all people,  
 $A(x, y) :=$  “ $x$  loves  $y$ ”,  $B(x) :=$  “ $x$  likes the Middle”, and  $\underline{c} :=$  “Raymond”.

(c) Consider the formula:  $(\forall x A(x)) \vee (\exists z B(f(z), \underline{c}))$ .

My model uses the universe of natural numbers  $\mathbb{N} = \{0, 1, 2, \dots\}$  and defines  $B(x, y) :=$  “ $x < y$ ”. Finish specifying my model in such a way that we prove the above statement is not logically valid.

**4. (25 points)** System K

(a) I have provided a proof of Lemma K17 below – except the justifications are missing.

Fill in the justifications.

**Lemma K17**     $\forall x \neg A(x) \vdash_K \neg \exists x A(x)$

1:  $\forall x \neg A(x)$

---

2:  $\forall x \neg A(x) \rightarrow \neg \neg \forall x \neg A(x)$

---

3:  $\neg \neg \forall x \neg A(x)$

---

4:  $\neg \exists x A(x)$

---

(b) Prove Lemma K22:  $\vdash_K \exists x \forall y A(x, y) \rightarrow \forall y \exists x A(x, y)$ .

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**1. (25 points)** “Fun” with Truth Tables!

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This statement is a    **Contingency**    /    **Contradiction**    /    **Tautology**.

- (b) Show that  $\neg a \wedge \neg b$  and  $\neg(a \vee b)$  are logically equivalent.

- (c) Consider the following statement:  $\neg b, a \rightarrow b \vdash \neg a$ . Circle the correct answers.

**IS**    /    **IS NOT**    a theorem of  $L$  by the    **Soundness**    /    **Completeness**    theorem.

## 2. (25 points) System L

(a) I have provided a proof of Lemma L8 below – except the justifications are missing.

Fill in the justifications and please **be specific!**

**Lemma L8**  $A \rightarrow B, B \rightarrow C \vdash_L A \rightarrow C$

1:  $A \rightarrow B$

---

2:  $B \rightarrow C$

---

3:  $A \rightarrow (B \rightarrow C)$

---

4:  $(A \rightarrow (B \rightarrow C)) \rightarrow ((A \rightarrow B) \rightarrow (A \rightarrow C))$

---

5:  $(A \rightarrow B) \rightarrow (A \rightarrow C)$

---

6:  $A \rightarrow C$

---

(b) Prove Lemma L10:  $\vdash_L (\neg B \rightarrow \neg A) \rightarrow (A \rightarrow B)$ .

[You may use the deduction theorem if you wish.]

*Hint/Suggestion:* Axiom 3 should be helpful.

**3. (25 points)** Models, Variables, Free for...

(a) Consider the formula:  $\exists x (A(f(z), x) \rightarrow \forall y (B(x, y) \wedge C(\underline{a}, y, w)))$ .

- i. Underline the scope of each quantifier in the above formula.
- ii. **Circle** all of the **bound variables** in the above formula.
- iii. Circle the correct answers and fill in the blanks below:

The formula above **IS** / **IS NOT** a sentence since \_\_\_\_\_.

The term  $t = g(y, \underline{a})$  **IS** / **IS NOT** free for  $z$  in the formula above.

(b) Consider the formula:  $(\exists x A(x, \underline{c})) \rightarrow (\forall y B(y))$ .

Translate the above formula into *plain English* when our model consists of the universe of all people,  
 $A(x, y) :=$  “ $x$  rescues  $y$ ”,  $B(x) :=$  “ $x$  will cheer”, and  $\underline{c} :=$  “Princess Fiona”.

(c) Consider the formula:  $(\forall x A(x, \underline{c})) \vee (\exists z B(f(z)))$ .

My model uses the universe of natural numbers  $\mathbb{N} = \{0, 1, 2, \dots\}$  and defines  $A(x, y) :=$  “ $x \geq y$ ”. Finish specifying my model in such a way that we prove the above statement is not logically valid.

**4. (25 points)** System K

(a) I have provided a proof of Lemma K17 below – except the justifications are missing.

Fill in the justifications.

**Lemma K17**     $\forall x \neg A(x) \vdash_K \neg \exists x A(x)$

1:  $\forall x \neg A(x)$

---

2:  $\forall x \neg A(x) \rightarrow \neg \neg \forall x \neg A(x)$

---

3:  $\neg \neg \forall x \neg A(x)$

---

4:  $\neg \exists x A(x)$

---

(b) Prove Lemma K22:  $\vdash_K \exists x \forall y A(x, y) \rightarrow \forall y \exists x A(x, y)$ .